Replica Management for the Grid

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Replica Management in Grids

- Need to manage large scientific computing datasets
  - Terabytes or petabytes shared by researchers around the world
  - Read-only data, "published" by experiments

- Replicate portions of the data set in multiple locations
  - Local control, reduce access times, provide fault tolerance

- Discover replicas and select the best replica for a necessary data transfer

Outline

- Data Intensive Applications: two examples

- Requirements for Data Grids

- A Replica Management System

  - Current implementation: Replica catalog and API for reliable replication

  - Address issues of replica location, file aggregation and reliable replication

- Replica Location Service

  - A flexible design framework

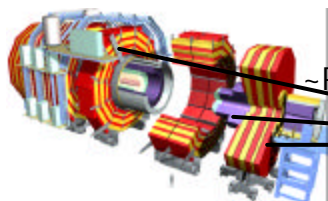 - Addresses issues of reliability, scalability and performance

Climate Modeling

- Detecting global climate change
- Simulate climate variability over long periods
  - Often 100 years
  - Long-duration computations: 1 month at 100 Gflops
  - Large output files: 10 terabytes

- Compare simulation results to observed variability

- With teraflop computers: requirements will increase by factor of 10 or more

- Need to publish, replicate and share these files with other researchers

Example: Data Replication for High Energy Physics

CMS

~PBytes/sec

Online System

~100 MBytes/sec

1 TIPS is approximately 25,000

SpecInt95 equivalents

There is a "bunch crossing" every 25 nsecs.

There are 100 "triggers" per second

Each triggered event is ~1 MByte in size

Offline Processor Farm

~20 TIPS

~100 MBytes/sec

Tier 0

CERN Computer Centre

HPSS

~622 Mbits/sec
or Air Freight (deprecated)

Tier 1

France Regional Centre   HPSS

Germany Regional Centre   HPSS

Italy Regional Centre   HPSS

FermiLab ~4 TIPS   HPSS

• • •

~622 Mbits/sec

Tier 2

Caltech ~1 TIPS   **Tier2 Centre ~1 TIPS**   **Centre TIPS**   **Centre TIPS**   **Centre TIPS**

~622 Mbits/sec

Institute ~0.25TIPS   tute   stitute   nstitute

Physics data cache

Physicists work on analysis "channels".

Each institute will have ~10 physicists working on one or more channels; data for these channels should be cached by the institute server

~1 MBytes/sec

Tier 4

Physicist workstations

Image courtesy Harvey Newman, Caltech   5

Data Grid Publication and Replication Requirements

- Terabytes or petabytes of data
  - Often read-only data, "published" by experiments

- Large data storage and computational resources shared by researchers around the world
  - Distinct administrative domains
  - Respect local and global policies governing how resources may be used

- Provide access to:
  - Raw experimental data
  - Simulation and analysis data products

the globus project
www.globus.org

Data Grid Requirements (Cont.)

- Management of data replication
  - Register and query physical copies of files
  - Reliably create and register new replicas
  - Select the best replica for a data transfer

- Security
  - Protect knowledge about existence of data

Fundamental Issues
for Replica Management

- *Location*: finding copies of files

- *Aggregation*: manage groups of files to improve convenience and scalability

- *Creation/Reliable replication*: copy files reliably and register them with management system

- *Scalability*: manage large numbers of files

- *Performance*: fast response time, large query and update rates

- *Reliability*: resilient to component failures

The Globus Architecture for Replica Management in Grids

- Managing multiple copies of data in wide area environments

- Identify <u>replica cataloging</u> and <u>reliable replication</u> as two fundamental services
  - Layer on other Grid services: GSI, transport, information service

- Used by higher-level services:
  - Replica selection
  - Automatic creation of new replicas

Our Data Model

- Data are organized into *files*

- Users group files into *collections*

- A *replica* or *location* is a subset of a collection stored on a particular physical storage system

- *Logical file name*: globally unique ID for a file within data grid's namespace

- *Physical file name*: location of an instance of a file on a particular storage system

- Maintain *mapping* between logical names for files and collections and one or more physical locations

The Replica Catalog

- Allows users to register replicas
- Answers queries about existing replicas

- Logical files
  - Entities with globally unique names, may have one or more physical instances

- Logical collection
  - Logical aggregations of groups of files (e.g., simulation timesteps)

- Location entries
  - All information required to map from logical to physical file names (hostname, port number, path…)
  - Location corresponds to one physical storage system

Replica Catalog Structure

the globus project
www.globus.org

```
                    ┌─────────────────────┐
                    │   Replica Catalog   │
                    └─────────────────────┘
```

Replica Catalog

Logical Collection

C02 measurements 1998

Logical Collection

C02 measurements 1999

Filename: Jan 1998
Filename: Feb 1998

…

Location
jupiter.isi.edu

Location
sprite.llnl.gov

**Logical
File Parent**

Filename: Mar 1998
Filename: Jun 1998
Filename: Oct 1998
Protocol: GridFTP
UrlConstructor:
GridFTP://jupiter.isi.edu/
    nfs/v6/climate

Filename: Jan 1998
…
Filename: Dec 1998
Protocol: ftp
UrlConstructor :
ftp://sprite.llnl.gov/
    pub/pcmdi

Logical File
Jan 1998

Size: 1468762

Logical File
Feb 1998

Reliable Replication
with the Replica Management API

- Reliable replication
  - Combines storage system operations with replica catalog updates
  - Create new replicas reliably
  - Automatically register them in the replica catalog

- Combined operations include:
  - Copy a file from one storage system to another and update the replica catalog
  - Delete a file from storage system and update catalog

- Reliability features
  - If can't complete operation, must **rollback** to previous consistent replica catalog state

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Replica Selection

Built on top of replica management

Given multiple physical copies of a desired file, want to select the "best" copy for a data transfer

- Select replica with best estimated performance

- Rely on *information services* that provide dynamic information about grid conditions
  - Storage system latency, bandwidth, load
  - Network latency and bandwidth
  - Authorization
  - User preferences

A Replica Location Service Framework

- Distribute replica management system to avoid single point of failure, performance bottleneck

- Applications may operate at different scales, have different resources and different tolerances to inconsistent RLS information

- We define a *flexible RLS framework*

- Allows users to make tradeoffs among:
  - Consistency, space overhead, reliability, update costs, query costs

- By different combinations of 5 essential elements, the framework supports a variety of RLS designs

Five Essential Elements of a Flexible RLS Framework

1. **Reliable Local State**

   Maintains consistent information about replicas at a single replica site

   – Updated when files are created or destroyed on local storage system

   – Contains mappings between LFNs & PFNs

   – Answers queries

2. **Global State with Relaxed Consistency**

   Implement as set of one or more Global Replica Index Nodes

   – Contain some LFN,replica site mappings

   – Accept periodic inputs from sites describing state

   – Answer queries for replicas associated with an LFN

Five Essential Elements of a Flexible RLS Framework

3. Soft State mechanisms for maintaining global state

   Soft state: information that times out and must be periodically refreshed

   – Stale information removed implicitly via timeouts
   – Index node state need not be persistent

4. Compression of State Updates
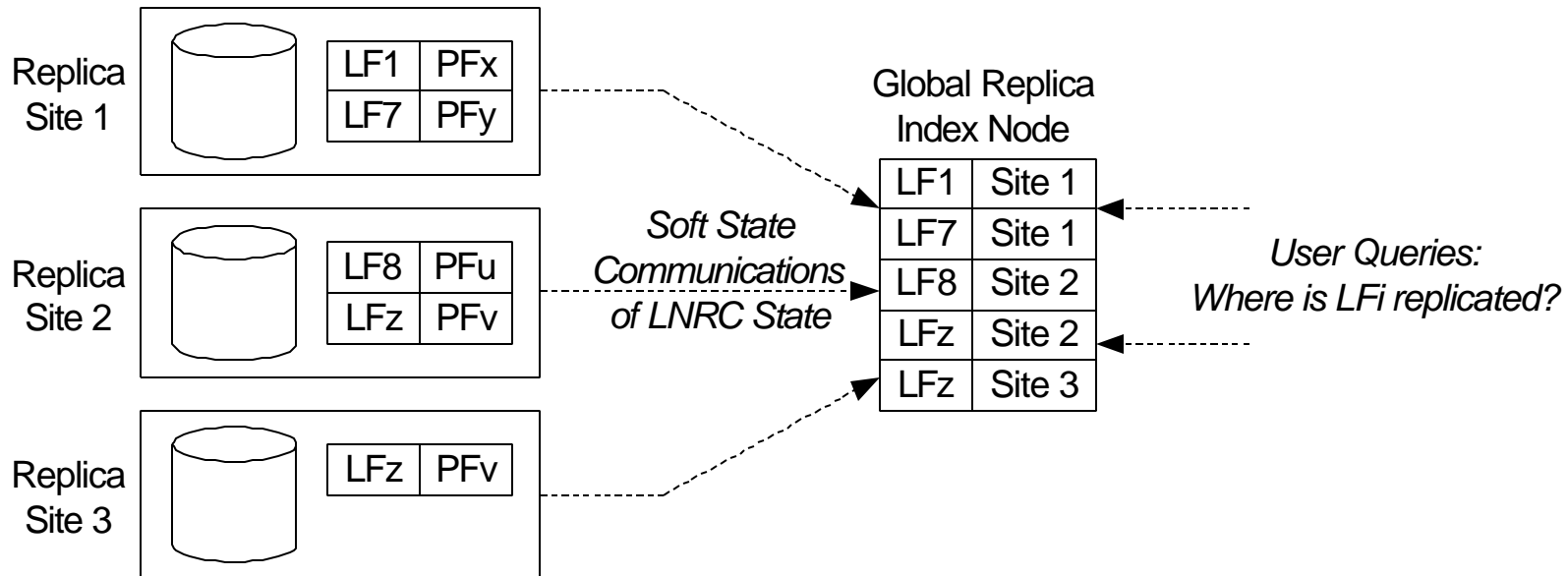
   Optional mechanism for reducing:

   – communication requirements for state updates
   – storage system requirements on GRINs

5. Membership Protocol

   For locating participating replica sites and GRINs

Example 1: A Centralized, Nonredundant Global Index

Replica Site 1

LF1	PFx
LF7	PFy

Replica Site 2

LF8	PFu
LFz	PFv

Replica Site 3

LFz	PFv

Soft State Communications of LNRC State

Global Replica Index Node

LF1	Site 1
LF7	Site 1
LF8	Site 2
LFz	Site 2
LFz	Site 3

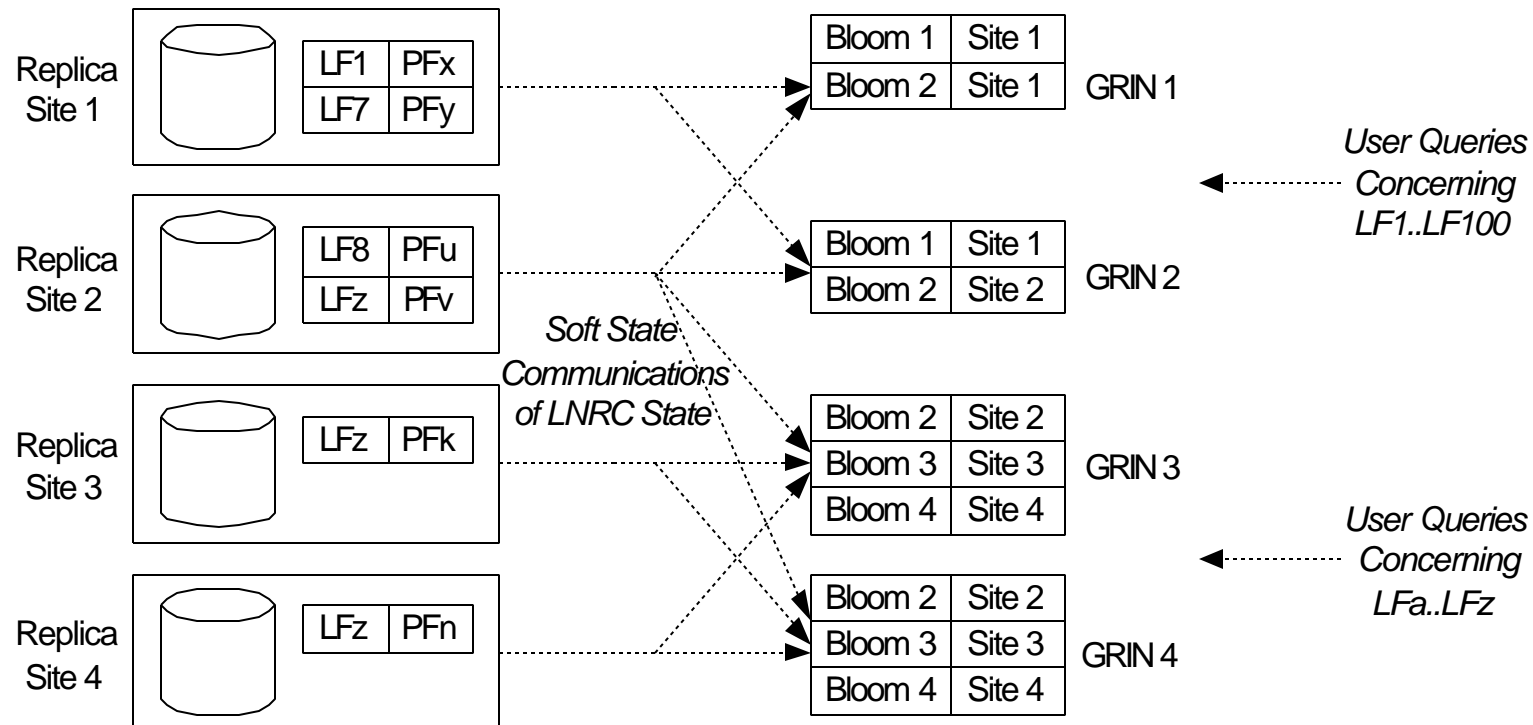*User Queries: Where is LFi replicated?*

All updates sent to a centralized GRIN

Not scalable: All queries serviced by a single index

Not reliable: Single point of failure

Example 2: An RLS with LFN Partitioning, Redundancy and Bloom Filter Compression

LF1	PFx
LF7	PFy

Replica Site 1

Bloom 1	Site 1
Bloom 2	Site 1

GRIN 1

LF8	PFu
LFz	PFv

Replica Site 2

Bloom 1	Site 1
Bloom 2	Site 2

GRIN 2

Soft State Communications of LNRC State

LFz	PFk

Replica Site 3

Bloom 2	Site 2
Bloom 3	Site 3
Bloom 4	Site 4

GRIN 3

LFz	PFn

Replica Site 4

Bloom 2	Site 2
Bloom 3	Site 3
Bloom 4	Site 4

GRIN 4

User Queries Concerning LF1..LF100

User Queries Concerning LFa..LFz

- Updates to specific, redundant GRINs based on LFN
- More scalable, reliable
- Limited storage and communication costs

Summary

- **Replica Management is a challenging problem for data-intensive applications**
  - Terabytes and petabytes of data
  - Replicated and shared by researchers around the world

- **Globus replica management**
  - Replica catalog
  - Reliable replication API

- **Replica Location Service**
  - Framework for flexible design of distributed replica location services
  - Reliable local state, relaxed global state, soft state updates, compression, membership